

Sudoku Players' Forums

Rating rules / Puzzles. Ordering the rules

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Author

Message

eleven

Di Posted: Wed Jul 08, 2009 2:53 pm Post subject:

(2) quote

denis_berthier wrote:

Joined: 10 Feb 2008 Posts: 325

But the version of SE I have (downloaded from gsf's page, with serate added) works fine and gives 7.1

Ah, thank you, it was fixed in the last version 1.2.1 (i had 1.2).

Changes wrote:

Sudoku Explainer: Changes since version 1.2

Added an additional check to fix false positives in BUG detection

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denis_berthier

□ Posted: Wed Jul 08, 2009 2:55 pm Post subject:



Joined: 19 Jun 2007

Posts: 673

Location: Paris, France

m_b_metcalf wrote:

denis_berthier wrote:

Mike:

- do you allow me to put your 64,000 bottom-up collection on my web pages? (with a reference to you, of course)
- is the description of your generator OK?

Denis, yes and yes.

Done.

I don't think I'll be able to consider larger collections, because my computing resources are limited and holidays are approaching.

Also, I think we wouldn't learn much more from the existing generators.

If a generator of a completely new kind appeared, it would be interesting to add it to our list and resume these comparisons.

Allan,

One of the big surprises I forgot to mention is about your generator. As the generation phase is very different from all the other generators, I thought it would lead to different results. But it is incredibly close to sudogen0_1M.

What seems the most important factor is not how the complete grids are (randomly) generated but rather the fact that we build complete grids before we start deleting redundant clues.

I nevertheless continue computations for a little more puzzles.

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ronk

Di Posted: Wed Jul 08, 2009 3:45 pm Post subject:

(Q) quote

*** off-topic ***

Joined: 02 Nov 2005 Posts: 2390

Location: Southeastern

USA

eleven wrote:

if you generated suexg14-0_1M with seed 0, and let it rate by Sudoku Explainer, keep an eye on puzzle nr. 737294. SE failed to rate it due to a bug in its BUG algorithm. The correct rating would be 7.1.

Code:

48	1	7	3	68	2	46	9	5	
5	2	3	7	9	46	46	8	1	
9	48	6	5	1	48	7	2	3	ļ
2	5	8	-+ 46	7	1	-+ 3	46	9	Ī
L	7	4	68	3	9	2	5	68	
5	3	9	2	48	5	18	14	7	ļ
- 3	#468	5	-+ 9	#468	7	1-8	16	2	1
18	68	2	1	5	3	9	7	#468	
7	9	1	48	2	68	5	3	#468	İ

It says, that in one of the marked 3-value cells there must be 8, otherwise there would be a BUG. But this is wrong.

ronk or someone else, can you explain, why the rule of thumb (2 more in the line/col/box) does not work here ?

After removing all the presumed extra candidates, every candidate must exist exactly twice (or not at all) in each row, column and box. For the above, if all four of the presumed extra 8s were removed, there would still be three 6s in each of r7, c9 and b9.

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eleven

□ Posted: Wed Jul 08, 2009 4:33 pm Post subject:



ronk wrote:

Joined: 10 Feb 2008 Posts: 325 *** off-topic ***

After removing all the presumed extra candidates, every candidate must exist exactly twice (or not at all) in each row, column and box. For the above, if all four of the presumed extra 8s were removed, there would still be three 6s in each of r7, c9 and b9.

Got it, thanks. I have to be careful, when BUG cells are in the same unit.

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StrmCkr

□ Posted: Wed Jul 08, 2009 6:53 pm Post subject:



my generator: written in free {turbo}pascal..

Joined: 05 Sep 2006

Posts: 473

Location: Winterpeg

uses the coded solving techniques to verify the grid is unique. and for deleting extra clues.. thats its downfall so far...

notes: setpm updates the grid. wpm writes it to screen. initiate resets everything back to zero.

wpm also calculates the number of sovled cells, and setpm calulates the #of pm's.

pm:[0..8,0..8,1..9] stores all the pencilmarks. {as x,y,N} s:[0..8,0..8,0] stores the active candidates.

solve functions reduces pm states.

Code:

procedure generate;
var

```
xn,yn,n,ns,temp,xn2,yn2:integer;
gen: array[0..8,0..8,0..0] of integer;
begin
for xn:=0 to 8 do
for yn:= 0 to 8 do
 gen[xn,yn,0]:=0;
initiate;
setpm;
wpm;
while (( count < 81) ) do
begin
randomize;
temp:=0;
repeat
  xn := (random(9));
  yn:= (random(9));
  for n:= 1 to 9 do
   if pm[xn,yn,N]>0
     then temp:=1+temp;
until (s[xn,yn,0]=0) and (temp>1);
 if s[xn,yn,0]=0
then
repeat
    n := (random(9));
until (pm[xn,yn,(N+1)]=(N+1));
s[xn,yn,0]:=(N+1);
 gen[xn,yn,0]:=(N+1);
setpm;
solve;
 setpm;
cal;
 if ((countpm=0) and (count <> 81))
 then
 begin
   initiate;
    for xn:=0 to 8 do
     for yn:=0 to 8 do
      gen[xn,yn,0]:=0;
    setpm;
    cal;
    end;
  for xn:=0 to 8 do
   for yn:= 0 to 8 do
     if (nm[xn,yn,1]=0) and (s[xn,yn,0]=0)
      then
        begin
          initiate;
         for xn2:=0 to 8 do
          for yn2:=0 to 8 do
           gen[xn2,yn2,0]:=0;
           setpm;
           cal;
          end;
end;
n := 0;
repeat
```

```
initiate;
for xn := 0 to 8 do
for yn:= 0 to 8 do
   s[xn,yn,0]:=(gen[xn,yn,0]);
cal:
ns:=count;
randomize;
  repeat
  xn:= (random(9));
  yn:= (random(9));
  until (s[xn,yn,0]>0);
  s[xn,yn,0]:=0;
  setpm:
   solve;
    setpm;
     cal;
  if count=81
  then
  begin
    gen[xn,yn,0]:=0;
   end
  else
  n := (1+N);
until (n=ns);
initiate;
for xn := 0 to 8 do
for yn:= 0 to 8 do
   s[xn,yn,0] := (gen[xn,yn,0]);
setpm;
wpm;
end;
```

Last edited by StrmCkr on Wed Jul 08, 2009 7:06 pm; edited 3 times in total

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Location: Southeastern



ronk

USA

□ Posted: Wed Jul 08, 2009 6:53 pm Post subject:



Allan Barker wrote:

Joined: 02 Nov 2005 Posts: 2390 denis_berthier wrote:

What would be your length for a Naked Pair?

Now I see what you mean. One answer is 2, because the definition of a naked pair is based on two cells that both contain a truth. The best answer is perhaps that the number of truths used in a logical deduction are what matter. In the above case, the NP has cleared the grid leaving behind a bi-local pair and only one is needed by the logic. Normally such an NP wouldn't be in the logic but if it were there for pratical reasons then the bi-local can be substituted. (Good example!).

Borrowing **Denis'** term, it's a limit case in "complementarity." The complement to an AHS(1) is an LS(2) in the above example. 9 Examples where the complement to an AHS(1) is LS(3) or LS(4), etc. are easy to find. As before, the **N** of LS(**N**) is the number of cells in the set.

[edit: To add scope to borrowing, quoted the term complementarity]

Last edited by ronk on Wed Jul 08, 2009 8:02 pm; edited 1 time in total

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denis_berthier

Di Posted: Wed Jul 08, 2009 7:33 pm Post subject:



Joined: 19 Jun 2007

Posts: 673

Location: Paris, France

ronk wrote:

Borrowing Denis' term, it's a limit case in complementarity. The complement to an AHS(1) is an LS(2) in the above example. $\stackrel{\text{LS}}{=}$ Examples where the complement to an AHS(1) is LS(3) or LS(4), etc. are easy to find. As before, the \mathbf{N} of LS(\mathbf{N}) is the number of cells in the set.

I doubt ever speaking of "a limit case in complementarity".

I've never spoken of complementarity for anything else than pure (i.e. not almost or almost-almost...) Naked or Hidden Subsets. For whips or braids (pure or with ALS), that's all we ever need.

I'm not even sure there is a consistent definition of complementarity for the almost-somethings.

But I'm almost certain one can give similar definitions for Allan's patterns, consistent with all the cases for which these patterns can be interpreted as whips(ALS) or braids(ALS). As I'm not very familiar with the technical details of his approach, I'll wait for his analysis of the question.

Anyway, the general rule that assigns a length to a pattern should be based on the general form of this pattern and not on "limit cases".

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denis_berthier

□ Posted: Thu Jul 09, 2009 4:24 am Post subject:



I think there's one more thing we could learn.

Joined: 19 Jun 2007

Posts: 673 Location: Paris, France The previous computations have shown that puzzles generated bottom-up tend to be easier than puzzles generated top-down.

This remains very mysterious.

The computations have also shown that bottom-up puzzles tend to have fewer clues; but, in spite of the small trend more clues -> harder, this is not enough to explain the mystery, because the above tendency is true for any fixed number of clues.

I imagined the following double generation process:

- 1) take the first part of a bottom-up generator and get a non complete grid G1 with a unique solution
- 2) with the second part of the generator obtain a minimal puzzle P1
- 3) let G2 be the complete (solution) grid obtained from G1
- 4) with exactly the same second part (this is very important) obtain a minimal puzzle P2

Question: is there any correlation between P1 and P2:

- number of clues
- SER
- NRCZT

Mike, do you think your bottom-up generator could be adapted to do this? (10,000 couples of puzzles would be enough for such computations.)

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StrmCkr

Di Posted: Thu Jul 09, 2009 6:25 am Post subject:



Quote:

Joined: 05 Sep 2006 Posts: 473 Location: Winterpeg

This remains very mysterious.

the mystery is the logic set used to solve the puzzle.

in top down builds you remove clues from a fixed point until logic sets are in place to produces a valid grid, which can mean harder move sets to.

in bottom up builds generator has less chance of creating complex logic sets until it hits a state that is unique.

my generator code has been producing a few 7+ er rated puzzles with few clues and some with many,

because it uses logic to reduce the pm states via logic before it places the next randomly generated digit.

which gives it a increased chance of randomly creating a harder puzzle.

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denis_berthier

D Posted: Thu Jul 09, 2009 6:35 am Post subject:



StrmCkr wrote:

Joined: 19 Jun 2007

Posts: 673 Location: Paris, France

Quote:

This remains very mysterious.

the mystery is the logic set used to solve the puzzle.

Not sure I understand this part of your post.

It seems clear to me that the only thing that counts for the solver part of any generator is its output:

- 0 (0 solution)
- 1 (1 solution)
- 2 (at least 2 solutions).

How this result is obtained doesn't matter.

Or am I missing anything? Are you suggesting that the solver part of some of the generators that have been discussed here don't produce the correct 0/1/2 value?

I understand though that the choice of the next clue may have an impact on the difficulty of the puzzles but aren't you thus introducing a strong bias?

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m_b_metcalf

Posts: 2191 Location: Berlin

Joined: 15 May 2006

Deposited: Thu Jul 09, 2009 6:45 am Post subject:



StrmCkr wrote: Quote:

This remains very mysterious.

the mystery is the logic set used to solve the puzzle.

My generators use no logic, rather backtracking preceded by singles solving, solely for efficiency reasons (to prune the search tree).

Regards,

Mike Metcalf

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StrmCkr

Posted: Thu Jul 09, 2009 6:53 am Post subject:



Quote:

Joined: 05 Sep 2006

Posts: 473

Location: Winterpeg

Or am I missing anything? Are you suggesting that the solver part of some of the generators that have been discussed here don't produce the correct 0/1/2 value?

no they solve the state counts correctly.

what i mean with a fixed point of space, there is already defined pattern logic in place or n many cover sets keeping it as unique or near unique depending on the starting fixed point.

reducing clues from that state allows more complex patterns to cover the cells.

if the pattern allows the puzzle to remain as a unique solution and no further clues can be reduced.

the grid will more likely have extremely complex logic validating the solution.

in an up build u have zero patterns in place for logic sets and must build the logic into a singular state.

building the logic upwards in a random fashion u have a reduced chance of having complex logic in place.

the total clue count would be a by produce of the logic sets in place ie the minimal needed given the sets used to create the unique state.

Quote:

My generators use no logic, rather backtracking preceded by singles solving, solely for efficiency reasons (to prune the search tree).

understood a recursive back track is far faster.

Ive been using logic sets to reduce the pm state, to only choose potions that are not covered yet.

the bias ism would be the limitation of implemented logic.

{my solver only has a limited amount ATM thus it cant fully optimize a gird.}

Last edited by StrmCkr on Thu Jul 09, 2009 7:02 am; edited 2 times in total

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denis_berthier

D Posted: Thu Jul 09, 2009 6:59 am Post subject:



StrmCkr,

Joined: 19 Jun 2007

Posts: 673

Location: Paris, France

I understand this may be interesting when the purpose is to find hard puzzles. But when the purpose is statistical classification and rating, as in this thread, one can't accept such a deliberate bias.

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StrmCkr

□ Posted: Thu Jul 09, 2009 7:04 am Post subject:



i think u missed my point,

Joined: 05 Sep 2006

Posts: 473 Location: Winterpeg i am attempting to answering this question.

Quote:

The previous computations have shown that puzzles generated bottom-up tend to be easier than puzzles generated top-down.

This remains very mysterious.

your listing statics on a group of random puzzles, top up or bottom up building.

and comparing rating and statical mean of clues on minimal puzzles between the different generators.

in a top down it has fixed logic in place already. so it has an increased chance of generating more complex logic.

bottom up builds the logic sets randomly.

